

The proximal mapping

<http://bicmr.pku.edu.cn/~wenzw/opt-2020-fall.html>

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Outline

- 1 closed function
- 2 Conjugate function
- 3 Proximal Mapping

Closed set

a set C is closed if it contains its boundary:

$$x^k \in C, \quad x^k \rightarrow \bar{x} \quad \implies \quad \bar{x} \in C$$

operations that preserve closedness

- the intersection of (finitely or infinitely many) closed sets is closed
- the union of a finite number of closed sets is closed
- inverse under linear mapping: $\{x \mid Ax \in C\}$ is closed if C is closed

Image under linear mapping

the image of a closed set under a linear mapping is not necessarily closed

example(C is closed, $AC = \{Ax \mid x \in C\}$ is open):

$$C = \{(x_1, x_2) \in \mathbf{R}_+^2 \mid x_1 x_2 \geq 1\}, \quad A = \begin{bmatrix} 1 & 0 \end{bmatrix}, \quad AC = \mathbf{R}_{++}$$

sufficient condition: AC is closed if

- C is closed and convex
- and C does not have a recession direction in the nullspace of A .
i.e.,

$$Ay = 0, \quad \hat{x} \in C, \quad \hat{x} + \alpha y \in C \quad \forall \alpha \geq 0 \quad \implies \quad y = 0$$

in particular, this holds for any A if C is bounded

Closed function

definition: a function is closed if its epigraph is a closed set or if for each $\alpha \in \mathbb{R}$ the sublevel set $\{x \in \text{dom}f \mid f(x) \leq \alpha\}$ is a closed set.

- If f is a continuous function and $\text{dom}f$ is closed, then f is closed.
- If f is a continuous function and $\text{dom}f$ is open, then f is closed iff it converges to ∞ along every sequence converging to a boundary point of $\text{dom}f$

examples

- $f(x) = -\log(1 - x^2)$ with $\text{dom}f = \{x \mid |x| < 1\}$
- $f(x) = x \log x$ with $\text{dom}f = \mathbf{R}_+$ and $f(0) = 0$
- indicator function of a closed set C : $f(x) = 0$ if $x \in C = \text{dom}f$

not closed

- $f(x) = x \log x$ with $\text{dom}f = \mathbf{R}_{++}$ or $\text{dom}f = \mathbf{R}_+$ and $f(0) = 1$
- indicator function of a set C if C is not closed

Properties

sublevel sets: f is closed if and only if all its sublevel sets are closed

minimum: if f is closed with bounded sublevel sets then it has a minimizer

Weierstrass

Suppose that the set $D \subset E$ (a finite dimensional vector space over R^n) is nonempty and closed, and that all sublevel sets of the continuous function $f : D \rightarrow R$ are bounded. Then f has a global minimizer.

common operations on convex functions that preserve closedness

- sum: $f + g$ is closed if f and g are closed (and $\text{dom } f \cap \text{dom } g \neq \emptyset$)
- composition with affine mapping: $f(Ax + b)$ is closed if f is closed
- supremum: $\sup_{\alpha} f_{\alpha}(x)$ is closed if each function f_{α} is closed

Weierstrass 定理

考虑优化问题

$$\begin{aligned} \min_{x \in \mathbb{R}^n} \quad & f(x), \\ \text{s.t.} \quad & x \in \mathcal{X}, \end{aligned} \tag{1}$$

其中 $\mathcal{X} \subseteq \mathbb{R}^n$ 为可行域. 考虑一个适当且闭的函数 $f: \mathcal{X} \rightarrow (-\infty, +\infty]$, 假设下面三个条件中任意一个成立:

- 1 $\text{dom}f \stackrel{\text{def}}{=} \{x \in \mathcal{X} : f(x) < +\infty\}$ 是有界的;
- 2 存在一个常数 $\bar{\gamma}$ 使得下水平集 $C_{\bar{\gamma}} \stackrel{\text{def}}{=} \{x \in \mathcal{X} : f(x) \leq \bar{\gamma}\}$ 是非空且有界的;
- 3 f 是强制的, 即对于任一满足 $\|x^k\| \rightarrow +\infty$ 的点列 $\{x^k\} \subset \mathcal{X}$, 都有

$$\lim_{k \rightarrow \infty} f(x^k) = +\infty,$$

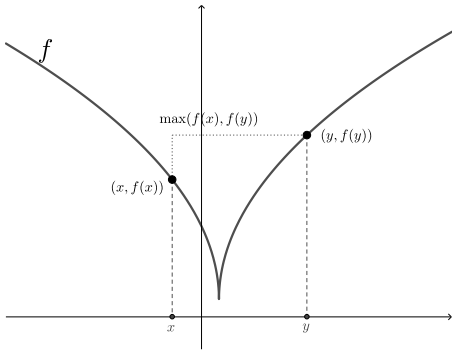
那么问题(1)的最小值点集 $\{x \in \mathcal{X} \mid f(x) \leq f(y), \forall y \in \mathcal{X}\}$ 非空且紧.

强拟凸函数

给定凸集 \mathcal{X} 和函数 $f: \mathcal{X} \rightarrow (-\infty, +\infty]$. 如果对任意的 $x \neq y$ 和 $\lambda \in (0, 1)$, 都有

$$f(\lambda x + (1 - \lambda)y) < \max\{f(x), f(y)\},$$

那么我们称函数 f 是强拟凸的.



- 任意强凸函数均为强拟凸的, 但凸函数并不一定是强拟凸的.

唯一性定理

唯一性定理: 对于问题(1), 设 \mathcal{X} 是 \mathbb{R}^n 的一个非空、紧且凸的子集, 如果 $f: \mathcal{X} \rightarrow (-\infty, +\infty]$ 是适当、闭且强拟凸函数, 那么存在唯一的 x^* 满足

$$f(x^*) < f(x), \quad \forall x \in \mathcal{X} \setminus \{x^*\}.$$

证明: 由Weierstrass定理知, 问题(1)至少存在一个全局极小解 x^* . 假设还有另外一个全局极小解 y^* , 那么 $f(x^*) = f(y^*)$. 根据强拟凸函数的定义, 对任意的 λ , 有

$$f(\lambda x^* + (1 - \lambda)y^*) < \max\{f(x^*), f(y^*)\} = f(x^*),$$

这与 x^* 的全局最优性矛盾.

- 对任何定义在有界凸集上的强凸函数 (如 $f(x) = x^2$), 其最优解都是唯一存在的.
- 一般凸函数最优解可能不唯一, 比如函数 $f(x) = \max\{x, 0\}$, 任意 $x \leq 0$ 都是 $f(x)$ 的最优解

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Conjugate function

the **conjugate** of a function f is

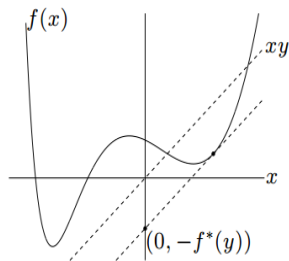
$$f^*(y) = \sup_{x \in \text{dom } f} (y^T x - f(x))$$

f^* is closed and convex even if f is not

Fenchel's s inequality

$$f(x) + f^*(y) \geq x^T y \quad \forall x, y$$

(extends inequality $x^T x/2 + y^T y/2 \geq x^T y$ to non-quadratic convex f)



Quadratic function

$$f(x) = \frac{1}{2}x^T A x + b^T x + c$$

strictly convex case ($A \succ 0$)

$$f^*(y) = \frac{1}{2}(y - b)^T A^{-1}(y - b) - c$$

general convex case ($A \succeq 0$)

$$f^*(y) = \frac{1}{2}(y - b)^T A^\dagger (y - b) - c, \quad \text{dom } f^* = \text{range}(A) + b$$

Negative entropy and negative logarithm

negative entropy

$$f(x) = \sum_{i=1}^n x_i \log x_i \quad f^*(y) = \sum_{i=1}^n e^{y_i - 1}$$

negative logarithm

$$f(x) = - \sum_{i=1}^n \log x_i \quad f^*(y) = - \sum_{i=1}^n \log(-y_i) - n$$

matrix logarithm

$$f(x) = - \log \det X \quad (\text{dom } f = \mathbf{S}_{++}^n) \quad f^*(Y) = - \log \det(-Y) - n$$

Indicator function and norm

indicator of convex set C : **conjugate** is support function of C

$$f(x) = \begin{cases} 0, & x \in C \\ +\infty, & x \notin C \end{cases} \quad f^*(y) = \sup_{x \in C} y^T x$$

norm: conjugate is indicator of unit dual norm ball

$$f(x) = \|x\| \quad f^*(y) = \begin{cases} 0, & \|y\|_* \leq 1 \\ +\infty, & \|y\|_* > 1 \end{cases}$$

(see next page)

proof: recall the definition of dual norm:

$$\|y\|_* = \sup_{\|x\| \leq 1} x^T y$$

to evaluate $f^*(y) = \sup_x (y^T x - \|x\|)$ we distinguish two cases

- if $\|y\|_* \leq 1$, then (by definition of dual norm)

$$y^T x \leq \|x\| \quad \forall x$$

and equality holds if $x = 0$; therefore $\sup_x (y^T x - \|x\|) = 0$

- if $\|y\|_* > 1$, there exists an x with $\|x\| \leq 1, x^T y > 1$; then

$$f^*(y) \geq y^T(tx) - \|tx\| = t(y^T x - \|x\|)$$

and *r.h.s.* goes to infinity if $t \rightarrow \infty$

The second conjugate

$$f^{**}(x) = \sup_{y \in \text{dom } f^*} (x^T y - f^*(y))$$

- $f^{**}(x)$ is closed and convex
- from Fenchel's inequality $x^T y - f^*(y) \leq f(x)$ for all y and x):

$$f^{**} \leq f(x) \quad \forall x$$

equivalently, $\text{epi } f \subseteq \text{epi } f^{**}$ (for any f)

- if f is closed and convex, then

$$f^{**}(x) = f(x) \quad \forall x$$

equivalently, $\text{epi } f = \text{epi } f^{**}$ (if f is closed convex); proof on next page

proof ($f^{**} = f$ if f is closed and convex): by contradiction
 suppose $(x, f^{**}(x)) \notin \text{epi } f$; then there is a strict separating
 hyperplane:

$$\begin{bmatrix} a \\ b \end{bmatrix}^T \begin{bmatrix} z - x \\ s - f^{**}(x) \end{bmatrix} \leq c \leq 0 \quad \forall (z, s) \in \text{epi } f$$

for some a, b, c with $b \leq 0$ ($b > 0$ gives a contradiction as $s \rightarrow \infty$)

- if $b < 0$, define $y = a/(-b)$ and maximize l.h.s. over $(z, s) \in \text{epi } f$:

$$f^*(y) - y^T x + f^{**}(x) \leq c/(-b) < 0$$

this contradicts **Fenchel's** inequality

- if $b = 0$, choose $\hat{y} \in \text{dom } f^*$ and add small multiple of $(\hat{y}, -1)$ to (a, b) :

$$\begin{bmatrix} a + \epsilon \hat{y} \\ -\epsilon \end{bmatrix}^T \begin{bmatrix} z - x \\ s - f^{**}(x) \end{bmatrix} \leq c + \epsilon(f^*(\hat{y}) - x^T \hat{y} + f^{**}(x)) < 0$$

now apply the argument for $b < 0$

Conjugates and subgradients

if f is closed and convex, then

$$y \in \partial f(x) \iff x \in \partial f^*(y) \iff x^T y = f(x) + f^*(y)$$

proof: if $y \in \partial f(x)$, then $f^*(y) = \sup_u (y^T u - f(u)) = y^T x - f(x)$

$$\begin{aligned} f^*(v) &= \sup_u v^T u - f(u) \\ &\geq v^T x - f(x) \\ &= x^T (v - y) - f(x) + y^T x \\ &= f^*(y) + x^T (v - y) \end{aligned} \tag{2}$$

for all v ; therefore, x is a subgradient of f^* at y ($x \in \partial f^*(y)$)

reverse implication $x \in \partial f^*(y) \Rightarrow y \in \partial f(x)$ follows from $f^{**} = f$

Some calculus rules

separable sum

$$f(x_1, x_2) = g(x_1) + h(x_2) \quad f^*(y_1, y_2) = g^*(y_1) + h^*(y_2)$$

scalar multiplication: (for $\alpha > 0$)

$$f(x) = \alpha g(x) \quad f^*(y) = \alpha g^*(y/\alpha)$$

addition to affine function

$$f(x) = g(x) + a^T x + b \quad f^*(y) = g^*(y - a) - b$$

infimal convolution

$$f(x) = \inf_{u+v=x} (g(u) + h(v)) \quad f^*(y) = g^*(y) + h^*(y)$$

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Proximal mapping

$$\text{prox}_f(x) = \underset{u}{\operatorname{argmin}} \left(f(u) + \frac{1}{2} \|u - x\|_2^2 \right)$$

if f is closed and convex then $\text{prox}_f(x)$ **exists** and is **unique** for all x

- existence: $f(u) + (1/2)\|u - x\|_2^2$ is closed with bounded sublevel sets
- uniqueness: $f(u) + (1/2)\|u - x\|_2^2$ is strictly (in fact, strongly) convex

subgradient characterization

$$u = \text{prox}_f(x) \iff x - u \in \partial f(u)$$

we are interested in functions f for which prox_f is inexpensive

Examples

quadratic function ($A \succeq 0$)

$$f(x) = \frac{1}{2}x^T Ax + b^T x + c, \quad \text{prox}_{tf}(x) = (I + tA)^{-1}(x - tb)$$

Euclidean norm: $f(x) = \|x\|_2$

$$\text{prox}_{tf}(x) = \begin{cases} (1 - t/\|x\|_2)x & \|x\|_2 \geq t \\ 0 & \textit{otherwise} \end{cases}$$

logarithmic barrier

$$f(x) = -\sum_{i=1}^n \log x_i, \quad \text{prox}_{tf}(x)_i = \frac{x_i + \sqrt{x_i^2 + 4t}}{2}, \quad i = 1, \dots, n$$

Some simple calculus rules

separable sum

$$f\left(\begin{bmatrix} x \\ y \end{bmatrix}\right) = g(x) + h(y), \quad \text{prox}_f\left(\begin{bmatrix} x \\ y \end{bmatrix}\right) = \begin{bmatrix} \text{prox}_g(x) \\ \text{prox}_h(y) \end{bmatrix}$$

scaling and translation of argument : with $\lambda \neq 0$

$$f(x) = g(\lambda x + a), \quad \text{prox}_f(x) = \frac{1}{\lambda}(\text{prox}_{\lambda^2 g}(\lambda x + a) - a)$$

scaling and translation of argument : with $\lambda > 0$

$$f(x) = \lambda g(x/\lambda), \quad \text{prox}_f = \lambda \text{prox}_{\lambda^{-1} g}(x/\lambda)$$

Addition to linear or quadratic function

linear function

$$f(x) = g(x) + a^T x, \quad \text{prox}_f = \text{prox}_g(x - a)$$

quadratic function: with $u > 0$

$$f(x) = g(x) + \frac{u}{2} \|x - a\|_2^2, \quad \text{prox}_f(x) = \text{prox}_{\theta g}(\theta x + (1 - \theta)a),$$

where $\theta = 1/(1 + u)$

Moreau decomposition

$$x = \text{prox}_f(x) + \text{prox}_{f^*}(x) \quad \forall x$$

- follows from properties of conjugates and subgradients:

$$\begin{aligned} u = \text{prox}_f(x) &\iff x - u \in \partial f(u) \\ &\iff u \in \partial f^*(x - u) \\ &\iff x - u = \text{prox}_{f^*}(x) \end{aligned}$$

- generalizes decomposition by orthogonal projection on subspaces:

$$x = P_L(x) + P_{L^\perp}(x)$$

if L is a subspace, L^\perp its orthogonal complement (this is Moreau decomposition with $f = I_L, f^* = I_{L^\perp}$)

Extended Moreau decomposition

for $\lambda > 0$

$$x = \text{prox}_{\lambda f}(x) + \lambda \text{prox}_{\lambda^{-1}f^*}(x/\lambda) \quad \forall x$$

proof: apply Moreau decomposition to λf

$$\begin{aligned} x &= \text{prox}_{\lambda f}(x) + \text{prox}_{(\lambda f)^*}(x) \\ &= \text{prox}_{\lambda f}(x) + \lambda \text{prox}_{\lambda^{-1}f^*}(x/\lambda) \end{aligned}$$

second line uses $(\lambda f)^*(y) = \lambda f^*(y/\lambda)$

Composition with affine mapping

- for general A , the prox-operator of

$$f(x) = g(Ax + b)$$

does not follow easily from the prox-operator of g

- however if $AA^T = (1/\alpha)I$, we have

$$\text{prox}_f(x) = (I - \alpha A^T A)x + \alpha A^T (\text{prox}_{\alpha^{-1}g}(Ax + b) - b)$$

example: $f(x_1, \dots, x_m) = g(x_1 + x_2 + \dots + x_m)$

$$\text{prox}_f(x_1, \dots, x_m)_i = x_i - \frac{1}{m} \left(\sum_{j=1}^m x_j - \text{prox}_{mg} \left(\sum_{j=1}^m x_j \right) \right)$$

proof: $u = \text{prox}_f(x)$ is the solution of the optimization problem

$$\begin{aligned} \min_{u,y} \quad & g(y) + \frac{1}{2} \|u - x\|_2^2 \\ \text{s.t.} \quad & Au + b = y \end{aligned}$$

with variables u, y

- eliminate u using the expression

$$\begin{aligned} u &= x + A^T(AA^T)^{-1}(y - b - Ax) \\ &= (I - \alpha A^T A)x + \alpha A^T(y - b) \end{aligned}$$

- optimal y is minimizer of

$$g(y) + \frac{\alpha^2}{2} \|A^T(y - b - Ax)\|_2^2 = g(y) + \frac{\alpha}{2} \|y - b - Ax\|_2^2$$

solution is $y = \text{prox}_{\alpha^{-1}g}(Ax + b)$

Projection on affine sets

hyperplane: $C = \{x | a^T x = b\}$ (with $a \neq 0$)

$$P_C(x) = x + \frac{b - a^T x}{\|a\|_2^2} a$$

affine set: $C = \{x | Ax = b\}$ (with $A \in \mathbb{R}^{p \times n}$ and $\text{rank}(A) = p$)

$$P_C(x) = x + A^T (AA^T)^{-1} (b - Ax)$$

inexpensive if $p \ll n$, or $AA^T = I, \dots$

Projection on simple polyhedral sets

halfspace: $C = \{x | a^T x \leq b\}$ (with $a \neq 0$)

$$P_C(x) = \begin{cases} x + \frac{b - a^T x}{\|a\|_2^2} a & \text{if } a^T x > b \\ x & \text{if } a^T x \leq b \end{cases}$$

rectangle: $C = [l, u] = \{l \preceq x \preceq u\}$

$$P_C(x)_i = \begin{cases} l_i & x_i \leq l_i \\ x_i & l_i \leq x_i \leq u_i \\ u_i & x_i \geq u_i \end{cases}$$

nonnegative orthant: $C = \mathbf{R}_+^n$

$$P_C(x) = x_+ \quad (x_+ \text{ is componentwise maximum of } 0 \text{ and } x)$$

probability simplex: $C = \{x | 1^T x = 1, x \geq 0\}$

$$P_C(x) = (x - \lambda 1)_+$$

where λ is the solution of the equation

$$1^T (x - \lambda 1)_+ = \sum_{i=1}^n \max\{0, x_k - \lambda\} = 1$$

probability simplex: $C = \{x | a^T x = b, l \leq x \leq u\}$

$$P_C(x) = P_{[l,u]}(x - \lambda a)$$

where λ is the solution of

$$a^T P_{[l,u]}(x - \lambda a) = b$$

Projection on norm balls

Euclidean ball: $C = \{x \mid \|x\|_2 \leq 1\}$

$$P_C(x) = \begin{cases} \frac{1}{\|x\|_2}x & \text{if } \|x\|_2 > 1 \\ x & \text{if } \|x\|_2 \leq 1 \end{cases}$$

1-norm ball: $C = \{x \mid \|x\|_1 \leq 1\}$

$$P_C(x)_k = \begin{cases} x_k - \lambda, & x_k > \lambda \\ 0, & -\lambda \leq x_k \leq \lambda \\ x_k + \lambda, & x_k < -\lambda \end{cases}$$

$\lambda = 0$ if $\|x\|_1 \leq 1$; otherwise λ is the solution of the equation

$$\sum_{k=1}^n \max\{|x_k| - \lambda, 0\} = 1$$

Projection on simple cones

second order cone $C = \{(x, t) \in \mathbf{R}^{n \times 1} \mid \|x\|_2 \leq t\}$

$$P_C(x, t) = (x, t) \quad \text{if } \|x\|_2 \leq t, \quad P_C(x, t) = (0, 0) \quad \text{if } \|x\|_2 \leq -t$$

and

$$P_C(x, t) = \frac{t + \|x\|_2}{2\|x\|_2} \begin{bmatrix} x \\ \|x\|_2 \end{bmatrix} \quad \text{if } \|x\|_2 > t, x \neq 0$$

positive semidefinite cone $C = \mathbf{S}_+^n$

$$P_C(X) = \sum_{i=1}^n \max\{0, \lambda_i\} q_i q_i^T$$

if $X = \sum_{i=1}^n \lambda_i q_i q_i^T$ is the eigenvalue decomposition of X

Support function

conjugate of support function of closed convex set is indicator function

$$f(x) = S_C(x) = \sup_{y \in C} x^T y, \quad f^*(y) = I_C(y)$$

prox-operator of support function: apply Moreau decomposition

$$\begin{aligned} \text{prox}_{tf} &= x - t \text{prox}_{t^{-1}f^*}(x/t) \\ &= x - tP_C(x/t) \end{aligned}$$

example: $f(x)$ is sum of largest r components of x

$$f(x) = x_{[1]} + \cdots + x_{[r]} = S_C(x), \quad C = \{y \mid 0 \leq y \leq 1, 1^T y = r\}$$

prox-operator of f is easily evaluated via projection on C

Norms

conjugate of norm is indicator function of dual norm ball:

$$f(x) = \|x\|, \quad f^*(y) = I_B(y) \quad (B = \{y \mid \|y\|_* \leq 1\})$$

prox-operator of norm: apply Moreau decomposition

$$\begin{aligned} \text{prox}_{tf} &= x - t \text{prox}_{t^{-1}f^*}(x/t) \\ &= x - tP_B(x/t) \\ &= x - P_{tB}(x) \end{aligned}$$

useful formula for $\text{prox}_{t\|\cdot\|}$ when projection on $tB = \{x \mid \|x\| \leq t\}$ is cheap

examples: $\|\cdot\|_1$, $\|\cdot\|_2$

Distance to a point

distance (in general norm)

$$f(x) = \|x - a\|$$

prox-operator: from page 23, with $g(x) = \|x\|$

$$\begin{aligned}\text{prox}_{tf} &= a + \text{prox}_{tg}(x - a) \\ &= a + x - a - tP_B\left(\frac{x - a}{t}\right) \\ &= x - P_{tB}(x - a)\end{aligned}$$

B is the unit ball for the dual norm $\|\cdot\|_*$

Euclidean distance to a set

Euclidean distance (to a closed convex set C)

$$d(x) = \inf_{y \in C} \|x - y\|_2$$

prox-operator of distance

$$\text{prox}_{td}(x) = \theta P_C(x) + (1 - \theta)x, \quad \theta = \begin{cases} t/d(x) & d(x) \geq t \\ 1 & \text{otherwise} \end{cases}$$

prox-operator of squared distance: $f(x) = d(x)^2/2$

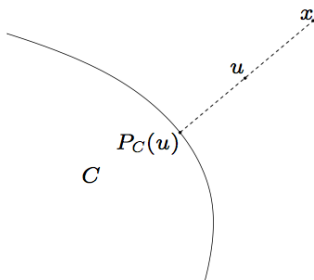
$$\text{prox}_{tf} = \frac{1}{1+t}x + \frac{t}{1+t}P_C(x)$$

proof (expression for $\text{prox}_{td}(x)$)

- if $u = \text{prox}_{td}(x) \notin C$, then from the definition and subgradient for d

$$x - u = \frac{t}{d(u)}(u - P_C(u))$$

implies $P_C(u) = P_C(x)$, $d(x) \geq t$, u is convex combination of x , $P_C(x)$



- if $u \in C$ minimizes $d(u) + (1/(2t))\|u - x\|_2^2$, then $u = P_C(x)$

proof (expression for $\text{prox}_{tf}(x)$ when $f(x) = d(x)^2/2$)

$$\begin{aligned}\text{prox}_{tf}(x) &= \arg \min_u \left(\frac{1}{2}d(u)^2 + \frac{1}{2t}\|u - x\|_2^2 \right) \\ &= \arg \min_u \inf_{v \in C} \left(\frac{1}{2}\|u - v\|_2^2 + \frac{1}{2t}\|u - x\|_2^2 \right)\end{aligned}$$

optimal u as a function of v is

$$u = \frac{t}{t+1}v + \frac{1}{t+1}x$$

optimal v minimizes

$$\frac{1}{2} \left\| \frac{t}{t+1}v + \frac{1}{t+1}x - v \right\|_2^2 + \frac{1}{2t} \left\| \frac{t}{t+1}v + \frac{1}{t+1}x - x \right\|_2^2 = \frac{1}{2(1+t)} \|v - x\|_2^2$$

Over C , i.e., $v = P_C(x)$

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