

Asymmetric Proximal Point Algorithms with Moving Proximal Centers

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- 2 APPA with Moving Proximal Centers
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- 6 Conclusion

$$(x - x^*)^T F(x^*) \geq 0, \quad \forall x \in \Omega, \quad (1)$$

- Ω : a nonempty, closed and convex set in \mathcal{R}^N ;
- F : a continuous and monotone mapping defined on \mathcal{R}^N .

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Examples: Equations, Complementarity Problems, Constrained Optimization Problems, Saddle Point Problems, etc.

Proximal Point Algorithms

- Classical PPA:

$$(x - x^{k+1})^T \left(F(x^{k+1}) + \frac{1}{c_k}(x^{k+1} - x^k) \right) \geq 0, \quad \forall x \in \Omega, \quad (2)$$

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- A General Version:

$$(x - x^{k+1})^T (F(x^{k+1}) + M_k(x^{k+1} - x^k)) \geq 0, \quad \forall x \in \Omega, \quad (3)$$

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$$M_k := 1/c_k M,$$

$$(x - x^{k+1})^T (c_k F(x^{k+1}) + M(x^{k+1} - x^k)) \geq 0, \quad \forall x \in \Omega. \quad (4)$$

Role of M in Algorithms:

Make the subproblems easier:

- Preconditioner: T. Pock and A. Chambolle [1].
- Decomposable of the subproblems.

[1] T. Pock and A. Chambolle. *Diagonal preconditioning for first order primal-dual algorithms in convex optimization*, IEEE Inter. Con. Comput. Vis., 2011, pp. 1762-1769.

A simple example

The saddle point problem:

$$\min_{u \in \mathcal{U}} \max_{v \in \mathcal{V}} \Phi(u, v) := f(u) + v^T A u - g(v). \quad (5)$$

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PDHG scheme:

$$\left\{ \begin{array}{l} \hat{u}^{k+1} = u^k - \tau A^T v^k, \\ u^{k+1} = \arg \min_{u \in \mathcal{U}} f(u) + \frac{1}{2\tau} \|u - \hat{u}^{k+1}\|^2, \\ \bar{u}^{k+1} = u^{k+1} + \theta(u^{k+1} - u^k), \\ \hat{v}^{k+1} = v^k + \sigma A \bar{u}^{k+1}, \\ v^{k+1} = \arg \min_{v \in \mathcal{V}} g(v) + \frac{1}{2\sigma} \|v - \hat{v}^{k+1}\|^2 \end{array} \right. \quad (6)$$

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PPA point of view

$$M := \begin{pmatrix} \frac{1}{\tau} I_m & -A^T \\ -\theta A & \frac{1}{\sigma} I_n \end{pmatrix}. \quad (7)$$

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- **Our Motivation:** Schemes without corrections.

Definitions

Let $F(\cdot)$ be a mapping from \mathcal{R}^N into \mathcal{R}^N . Then, $F(\cdot)$ is said to be

- Monotone: if

$$(F(x) - F(y))^T(x - y) \geq 0 \quad \forall x, y \in \mathcal{R}^N;$$

- Strongly monotone with modulus μ if

$$(F(x) - F(y))^T(x - y) \geq \mu \|x - y\|^2 \quad \forall x, y \in \mathcal{R}^N.$$

- Lipschitz continuous: if

$$\|F(x) - F(y)\| \leq L \|x - y\| \quad \forall x, y \in \mathcal{R}^N;$$

- Co-coercive with modulus $\sigma > 0$ if

$$(F(x) - F(y))^T(x - y) \geq \sigma \|F(x) - F(y)\|^2 \quad \forall x, y \in \mathcal{R}^N. \quad (8)$$

Reformulation

$M_s \equiv \frac{1}{2}(M + M^T)$: The symmetric part of M . Define

- $C := M_s^{-1/2}(M - M_s)M_s^{-1/2}$;
- $K \equiv M_s^{1/2}\Omega = \{M_s^{1/2}x : x \in \Omega\}$;
- $K^* \equiv M_s^{1/2}\Omega^* = \{M_s^{1/2}x^* : x^* \in \Omega^*\}$;
- $\tilde{F}_M(y) \equiv M_s^{-1/2}F(M_s^{-1/2}y)$, $\forall y \in K$;
- For any $x \in \Omega$, we define $y := M_s^{1/2}x$. Thus,

$$y' \equiv M_s^{1/2}x', \quad y^{k+1} \equiv M_s^{1/2}x^{k+1}, \quad \text{and} \quad y^k \equiv M_s^{1/2}x^k. \quad (9)$$

Then, a vector x^* is a solution of the variational inequality (1) if and only if $y^* := M_s^{1/2}x^*$ is a solution of the variational inequality of finding $y \in K$ such that

$$(y' - y^*)^T \tilde{F}_M(y^*) \geq 0, \quad \forall y' \in K. \quad (10)$$

The solution set of (10) is thus given by K^* .

A Useful Lemma

Lemma 1

Let the mapping $G : \mathcal{R}^n \rightarrow \mathcal{R}^n$ be co-coercive on a nonempty, closed, convex subset W in \mathcal{R}^n with modulus $\sigma > 1/2$, Then, for any $x, y, z \in W$, we have

$$(x - y)^T (G(z) - G(y)) \geq -\frac{\nu}{2} \|x - z\|^2, \quad (11)$$

where ν is an arbitrary number satisfying

$$0 < \frac{1}{2\sigma} < \nu < 1. \quad (12)$$

Lemma 2

[Proposition 12.5.3](Facchinei and Pang 2003) The mapping $\alpha\tilde{F}_M - C$ is co-coercive over K with modulus greater than $1/2$ if either one of the following two conditions holds:

- 1 F is Lipschitz continuous on Ω with modulus L and a $\tau \in (0, 1)$ exists such that, for all $y_1, y_2 \in K$,

$$\|\alpha\tilde{F}_M(y_2) - \alpha\tilde{F}_M(y_1) - M_s^{-1/2}MM_s^{-1/2}(y_2 - y_1)\| \leq \tau\|y_2 - y_1\|;$$

- 2 $F(x) = Qx + q$, for some positive semidefinite matrix $Q \equiv D + E$ with D positive definite and E symmetric, and $0 < \|I_n + H\| < 2$,

The Exact Version

Algorithm 1: APPA-MPC — The Exact Version

Step 0. Given a positive definite matrix $M \in \mathbb{R}^{n \times n}$, $\epsilon > 0$, and $x^0 \in \Omega$. Choose α such that $\alpha \tilde{F}_M - C$ is co-coercive over K with the modulus $\sigma > 1/2$. Set $k := 0$.

Step 1 [APPA Step]. Find the next iterate x^{k+1} by solving the subproblem

$$(x' - x^{k+1})^T [F(x^{k+1}) + M(x^{k+1} - (x^k - \alpha M^{-1} F(x^k)))] \geq 0, \forall x' \in \Omega. \quad (3.1)$$

Step 2 [Convergence Verification]. If

$$\|x^{k+1} - x^k\| \leq \epsilon, \quad (3.2)$$

then stop; Otherwise, set $k := k + 1$ and go to Step 1.

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Remark: The exact version of APPA-MPC ("APPA-MPC-E" in short):

The proximal center in x^k in classical PPA is shifted to $x^k - \alpha M^{-1} F(x^k)$ in (3.1).

Lemma 3

The subproblem (3.1) of the APPA-MPC-E can be rewritten as

$$(y' - y^{k+1})^T [\tilde{F}_M(y^{k+1}) + (I + C)(y^{k+1} - y^k) + \alpha \tilde{F}_M(y^k)] \geq 0, \forall y' \in K. \quad (13)$$

Convergence

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Lemma 4

Let y^ be an arbitrary solution point in K^* . Then, we have*

$$\|y^{k+1} - y^*\|^2 \leq \|y^k - y^*\|^2 - (1 - \nu) \|y^k - y^{k+1}\|^2, \quad (14)$$

where ν is an arbitrary number satisfying (12).

The Inexact Version

Algorithm 2: APPA-MPC — An Inexact Version

Step 0. Given a positive definite matrix $M \in \mathbb{R}^{n \times n}$, $\epsilon > 0$ and $x^0 \in \Omega$. Choose α such that $\alpha \tilde{F}_M - C$ is co-coercive over K with the modulus $\sigma > 1/2$. Choose a sequence of $\{\epsilon_k\}$ with $\epsilon_k > 0$ for all k and $\sum_{k=1}^{\infty} \epsilon_k < +\infty$. Set $k := 0$.

Step 1 [PPA Step]. Solve the subproblem (3.1) iteratively and find x^{k+1} such that

$$\max\{\|x^{k+1} - \bar{x}^{k+1}\|, \|F(x^{k+1}) - F(\bar{x}^{k+1})\|\} \leq \epsilon_k, \quad (4.3)$$

where

$$\bar{x}^{k+1} := P_{\Omega}[x^{k+1} - (F(x^{k+1}) + M(x^{k+1} - x^k) + \alpha F(x^k))]. \quad (4.4)$$

Step 2 [Convergence Verification]. If

$$\max\{\|x^{k+1} - x^k\|, \|x^{k+1} - \bar{x}^{k+1}\|\} \leq \epsilon, \quad (4.5)$$

then stop; Otherwise, set $k := k + 1$ and go to Step 1.

Convergence

Lemma 5

Let $\{x^k\}$ be the sequence generated by APPA-MPC-I. Then, there exist a positive scalar Γ such that the following inequality holds for any $y' \in K$:

$$(y' - \bar{y}^{k+1})^T (\tilde{F}_M(\bar{y}^{k+1}) + (I + C)(\bar{y}^{k+1} - y^k) + \alpha \tilde{F}_M(y^k)) \geq -\epsilon_k \|y' - \bar{y}^{k+1}\|^2 - \Gamma \epsilon_k. \quad (15)$$

Convergence

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Theorem 6

The sequence $\{x^k\}$ generated by the APPA-MPC-I globally converges to a solution point of the variational inequality (1).

Theorem 7

Let $\{x^k\}$ be the sequence generated by the APPA-MPC-E. For an integer $t > 0$, let

$$y_t := \frac{1}{t+1} \sum_{k=0}^t y^{k+1}, \quad (16)$$

then we have $y_t \in K$ and

$$(y_t - y')^T \tilde{F}_M(y') \leq \frac{1}{2(1+\alpha)(t+1)} \|y' - y^0\|^2, \quad \forall y' \in K. \quad (17)$$

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Lemma 8 (Theorem 2.3.5, Facchinei and Pang (2003))

The solution set K^* of the variational inequality is convex and it can be characterized as

$$K^* = \bigcap_{y' \in K} \{y^* \in K \mid (y' - y^*)^T \tilde{F}_M(y') \geq 0.\} \quad (18)$$

Theorem 9

Let $\{x^k\}$ be the sequence generated by the APPA-MPC-I. For an integer $t > 0$, let

$$y_t := \frac{1}{t+1} \sum_{k=0}^t \bar{y}^{k+1}, \quad (19)$$

then we have $y_t \in K$ and

$$(y_t - y')^T \tilde{F}_M(y') \leq \frac{1}{2(1+\alpha)(t+1)} \left\{ \|y' - y^0\|^2 + N(y') \sum_{k=0}^t \epsilon_k \right\}, \quad \forall y' \in K.$$

Theorem 10

Suppose that F is strongly monotone with the modulus $\mu_f > 0$ and Lipschitz continuous with the constant L_f . Let $\{x^k\}$ be the sequence generated by the proposed APPA-MPC-E. Then it converges to a solution point in Ω^ linearly.*

Theorem 10

Suppose that F is strongly monotone with the modulus $\mu_f > 0$ and Lipschitz continuous with the constant L_f . Let $\{x^k\}$ be the sequence generated by the proposed APPA-MPC-E. Then it converges to a solution point in Ω^ linearly.*

Theorem 11

Suppose that F is strongly monotone with the modulus $\mu_f > 0$ and Lipschitz continuous with the constant L_f . Moreover, assume that

$$\epsilon_k \leq \frac{3(1-\nu)}{8\Gamma+2} \|y^k - \bar{y}^{k+1}\|^2, \quad \forall k \geq 0.$$

Let $\{x^k\}$ be the sequence generated by the proposed APPA-MPC-I. Then it converges to a solution point in Ω^ linearly.*

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The saddle-point problem

Recall that the problem is

$$\min_{u \in \mathcal{U}} \max_{v \in \mathcal{V}} \Phi(u, v) := f(u) + v^T A u - g(v), \quad (20)$$

where $f : \mathcal{U} \rightarrow \mathcal{R}$, $g : \mathcal{V} \rightarrow \mathcal{R}$ are convex but not necessarily smooth functions; $\mathcal{U} \subseteq \mathcal{R}^n$ and $\mathcal{V} \subseteq \mathcal{R}^m$ are two nonempty, closed and convex sets; and $A \in \mathcal{R}^{m \times n}$.

The algorithm

Algorithm 3: A Primal-Dual Hybrid Gradient Algorithm for the Saddle Point Problem (1.7)

Input: Choose $u^0 \in \mathbb{R}^n$, $v^0 \in \mathbb{R}^m$, $s_u^0 \in \partial f(u^0) + N_{\mathcal{U}}(u^0)$, $s_v^0 \in \partial g(v^0) + N_{\mathcal{V}}(v^0)$, $\theta \in \mathbb{R}$, and $\tau, \sigma > 0$ such that the matrix M defined by (1.9) is positive definite. Choose α such that $\alpha \tilde{F}_M - C$ is co-coercive over K with the modulus $\sigma > 1/2$.

```
1 while Not converged do
2    $\tilde{u}^{k+1} = u^k - \tau A^T v^k - \tau \alpha (s_u^k + A^T v^k)$ 
3    $u^{k+1} = \arg \min_{u \in \mathcal{U}} f(u) + \frac{1}{2\tau} \|u - \tilde{u}^{k+1}\|^2$ 
4    $s_u^{k+1} = \frac{1}{\tau} (\tilde{u}^{k+1} - u^{k+1})$ 
5    $\bar{u}^{k+1} = u^{k+1} + \theta (u^{k+1} - u^k)$ 
6    $\tilde{v}^{k+1} = v^k + \sigma A \bar{u}^{k+1} - \sigma \alpha (s_v^k - A \bar{u}^k)$ 
7    $v^{k+1} = \arg \min_{v \in \mathcal{V}} g(v) + \frac{1}{2\sigma} \|v - \tilde{v}^{k+1}\|^2$ 
8    $s_v^{k+1} = \frac{1}{\sigma} (\tilde{v}^{k+1} - v^{k+1})$ 
9 end
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The Problems

$$\min \left\{ \sum_{i=1}^m \theta_i(x_i) \mid \sum_{i=1}^m A_i x_i = b, x_i \in \mathcal{X}_i; i = 1, \dots, m \right\}, \quad (21)$$

where $\theta_i : \mathcal{R}^{n_i} \rightarrow \mathcal{R}$ are closed proper convex functions (not necessarily smooth); $A_i \in \mathcal{R}^{l \times n_i}$; $\mathcal{X}_i \subset \mathcal{R}^{n_i}$ are closed and convex nonempty sets; $b \in \mathcal{R}^l$; and $\sum_{i=1}^m n_i = n$.

The Decomposition Algorithm

$$M = \begin{pmatrix} \kappa I & -\beta A_1^T A_2 & \cdots & -\beta A_1^T A_m & 0 \\ 0 & \kappa I & \cdots & -\beta A_2^T A_m & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \kappa I & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{\beta} I \end{pmatrix}. \quad (22)$$

The algorithm

Algorithm 4: A Splitting Algorithm for (7.6)

Input: Choose $x_i^0 \in \mathbb{R}^{n_i}$, $s_i^0 \in \partial\theta_i(x_i^0)$, $i = 1, \dots, m$, and $\beta > 0$. Choose $\kappa > 0$ such that the matrix M defined by (7.12) is positive definite. Choose α such that $\alpha\tilde{F}_M - C$ is co-coercive over K with the modulus $\sigma > 1/2$.

```
1 while Not converged do
2   for  $i = 1, \dots, m$  do
3      $\tilde{x}_i^{k+1} := x_i^k - \frac{1}{\kappa}\alpha(s_i^k - A_i^T\lambda^k)$ 
4      $\phi_i^k := \sum_{j=1}^{i-1} A_j x_j^{k+1} + \sum_{j=i+1}^m A_j x_j^k - b - \frac{1}{\beta}\lambda^k$ 
5      $x_i^{k+1} = \arg \min_{x_i \in \mathcal{X}_i} \theta_i(x_i) + \frac{\beta}{2}\|A_i x_i + \phi_i^k\|^2 + \frac{\kappa}{2}\|x_i - \tilde{x}_i^{k+1}\|^2$ 
6      $\lambda^{k+1} = \lambda^k - \beta(\sum_{i=1}^m A_i x_i^{k+1} - b)$ 
7      $s_i^{k+1} = \kappa(\tilde{x}_i^{k+1} - x_i^{k+1}) - \beta A_i^T(A_i x_i^{k+1} + \phi_i^k)$ 
8   end
9 end
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Thank you !