

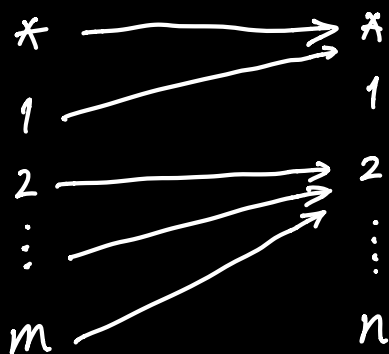
III. Derived Rings

1. ∞ -Operads

Our language to make sense of commutativity and associativity up to homotopy.

Def: Segal's category $\mathcal{F}in_*$ of pointed finite sets:

$$\begin{cases} \text{Obj: } \langle n \rangle = \{*, 1, \dots, n\}, & n \geq 0 \\ \text{Mor: } \alpha: \langle m \rangle \rightarrow \langle n \rangle & \text{s.t. } \alpha(*) = * \end{cases}$$



$\forall 1 \leq i \leq n$, let $\rho^i: \langle n \rangle \rightarrow \langle 1 \rangle$ sending $i \mapsto 1$, others $\mapsto *$.

A mor $f: \langle m \rangle \rightarrow \langle n \rangle$ in $\mathcal{F}in_*$ is **inert** if $\forall i \in \langle n \rangle^\circ = \{1, \dots, n\}$, $f^{-1}\{i\}$ has exactly one element.

Def: An ∞ -operad is a functor $p: \mathcal{O}^\otimes \rightarrow \mathcal{F}in_*$ between ∞ -cats, s.t.

(1) \forall inert mor $f: \langle m \rangle \rightarrow \langle n \rangle$ in $\mathcal{F}in_*$ \forall obj $C \in \mathcal{O}_{\langle m \rangle}^\otimes$, \exists a p -cocartesian mor $\bar{f}: C \rightarrow C'$ in \mathcal{O}^\otimes lifting f . Any such \bar{f} is called **inert**. In particular, f induces a functor $f_!: \mathcal{O}_{\langle m \rangle}^\otimes \rightarrow \mathcal{O}_{\langle n \rangle}^\otimes$

(2) Let $C \in \mathcal{O}_{\langle m \rangle}^{\otimes}$, $C' \in \mathcal{O}_{\langle n \rangle}^{\otimes}$, $f: \langle m \rangle \rightarrow \langle n \rangle$ in Fin_* .

Let $\text{Map}_{\mathcal{O}^{\otimes}}^f(C, C') \subset \text{Map}_{\mathcal{O}^{\otimes}}(C, C')$ be union of conn comp lying over f

Choose p -cocartesian morphisms $C' \rightarrow C'_i$ lying over the inert morphisms $p^i: \langle n \rangle \rightarrow \langle 1 \rangle$ for $1 \leq i \leq n$. Then the induced map

$$\text{Map}_{\mathcal{O}^{\otimes}}^f(C, C') \longrightarrow \prod_{1 \leq i \leq n} \text{Map}_{\mathcal{O}^{\otimes}}^{p^i \circ f}(C, C'_i)$$

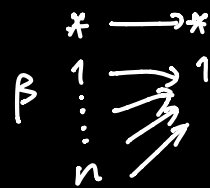
is a homotopy equivalence.

(3) $\forall n \geq 0$, the functors $\{p^i: \mathcal{O}_{\langle n \rangle}^{\otimes} \rightarrow \mathcal{O}_{\langle 1 \rangle}^{\otimes}\}_{1 \leq i \leq n}$ determine an equiv of ∞ -cats $\phi: \mathcal{O}_{\langle n \rangle}^{\otimes} \rightarrow \mathcal{O}^n$.

Notation: $\mathcal{O} := \mathcal{O}_{\langle 1 \rangle}^{\otimes}$ is called the **underlying ∞ -category of \mathcal{O}^{\otimes}** .

Given $X_1, \dots, X_n, Y \in \mathcal{O}$, let $\text{Mul}_{\mathcal{O}}(\{X_i\}_{1 \leq i \leq n}, Y)$ denote the union of components of $\text{Map}_{\mathcal{O}^{\otimes}}(\underbrace{X_1 \oplus \dots \oplus X_n}_{\mathcal{O}_{\langle n \rangle}^{\otimes}}, Y)$ which lies over the unique

mor $\beta: \langle n \rangle \rightarrow \langle 1 \rangle$ st. $\beta^{-1}\{*\} = \{*\}$.

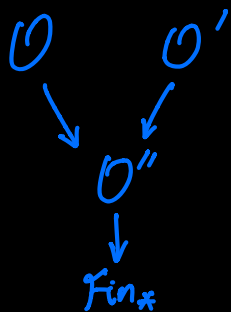


Idea: We think of an ∞ -operad as a category \mathcal{O} together with "multi-morphisms" $\text{Mul}_{\mathcal{O}}(\{X_i\}_{1 \leq i \leq n}, Y)$ whose compositions are associative up to homotopy.

Example: The **commutative ∞ -operad $\text{Comm}^{\otimes} := \text{Fin}_*$**

Def: $\mathcal{O}^{\otimes}, \mathcal{O}'^{\otimes}$ ∞ -operads. An ∞ -operad map from \mathcal{O}^{\otimes} to \mathcal{O}'^{\otimes} is a functor $f: \mathcal{O}^{\otimes} \rightarrow \mathcal{O}'^{\otimes}$ of ∞ -cats over Fin_* which carries inert mor in \mathcal{O}^{\otimes} to inert mor in \mathcal{O}'^{\otimes} .

Let $\text{Alg}_{\mathcal{O}}(\mathcal{O}')$ denote the full subcat of $\text{Fun}_{\text{Fin}_*}(\mathcal{O}^{\otimes}, \mathcal{O}'^{\otimes})$ spanned by ∞ -operad maps.



$$\text{Alg}_{\mathcal{O}/\mathcal{O}''}(\mathcal{O}')$$

$$\text{Alg}_{\mathcal{O}}(\mathcal{O}') := \text{Alg}_{\mathcal{O}/\mathcal{O}}(\mathcal{O}') \text{ where } \mathcal{O} = \mathcal{O}''.$$

Def: \mathcal{O}^{\otimes} ∞ -operad. A cocartesian fibration $p: \mathcal{C}^{\otimes} \rightarrow \mathcal{O}^{\otimes}$ is called a **cocartesian fibration of ∞ -operads** if the composition $\mathcal{C}^{\otimes} \rightarrow \mathcal{O}^{\otimes} \rightarrow \text{Fin}_*$ is an ∞ -operad. In this case, \mathcal{C} is called an **\mathcal{O} -monoidal ∞ -category**.

Idea: $\forall f \in \text{Mul}_{\mathcal{O}}(\{X_i\}_{1 \leq i \leq n}, Y)$, the cocartesian fibration p determines a functor $\otimes_f: \prod_{1 \leq i \leq n} \mathcal{C}_{X_i} \rightarrow \mathcal{C}_Y$

↑ fiber of \mathcal{C} at X_i

Example: A **symmetric monoidal ∞ -cat** \mathcal{C} is a Fin_* -monoid ∞ -cat, i.e. a cocartesian fibration $\mathcal{C}^{\otimes} \rightarrow \text{Fin}_*$ s.t. $\forall n \geq 0$, the maps

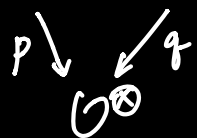
$\{\rho_i: \langle n \rangle \rightarrow \langle 1 \rangle\}_{1 \leq i \leq n}$ induces functors $\rho_i^{\otimes}: \mathcal{C}_{\langle n \rangle}^{\otimes} \rightarrow \mathcal{C}_{\langle 1 \rangle}^{\otimes}$ which determine an equivalence $\mathcal{C}_{\langle n \rangle}^{\otimes} \simeq (\mathcal{C}_{\langle 1 \rangle}^{\otimes})^n$.

The morphisms $\alpha: \langle 0 \rangle \rightarrow \langle 1 \rangle$ and $\beta: \langle 2 \rangle \rightarrow \langle 1 \rangle$ in Fin_* determines $(\text{Alg}(\mathcal{C}) := \text{Alg}_{\text{Fin}_*}(\mathcal{C}))$ **∞ -category of commutative algebra objects of \mathcal{C}** .

functors $\Delta^{\circ} \rightarrow \mathcal{C}$ $\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$

Corresponding to the unit object $1 \in \mathcal{C}$ and the tensor product.

Def: \mathcal{C}^{\otimes} \mathcal{D}^{\otimes} \mathcal{O} -monoidal ∞ -categories



$f \in \text{Alg}_{\mathcal{C}/\mathcal{O}}(\mathcal{D})$ is a \mathcal{O} -monoidal functor if it carries p -cocartesian mor to q -cocartesian mor.

In the special case $\mathcal{O}^{\otimes} = \text{Fin}_*$, we say symmetric monoidal functor.

Def: \mathcal{C} ∞ -category. A symmetric monoidal structure on \mathcal{C} is cartesian

if • the unit obj $1_{\mathcal{C}} \in \mathcal{C}$ is final

• $\forall C, D \in \mathcal{C}$, the canonical maps $C \simeq C \otimes 1_{\mathcal{C}} \leftarrow C \otimes D \rightarrow 1_{\mathcal{C}} \otimes D \simeq D$

exhibit $C \otimes D$ as a product of C and D in the ∞ -cat \mathcal{C} .

Dually \rightsquigarrow cocartesian symmetric monoidal structure.

HA2.4.1: \mathcal{C} ∞ -cat admitting finite products. Then \mathcal{C} admits a cartesian symmetric monoidal structure, unique up to equivalence.

In this case, the ∞ -cat $\mathcal{C}(\text{Alg}(\mathcal{C}))$ admits a direct description in terms of monoids.

2. Algebras and modules

\mathcal{C}^{\otimes} symmetric monoidal ∞ -cat $\mathcal{O} := \text{Fin}_*$

$$\text{Alg}(\mathcal{C}) := \text{Alg}_{\mathcal{O}}(\mathcal{C})$$

Evaluate at $\langle 1 \rangle \in \text{Fin}_* \rightsquigarrow$ Forgetful functor $\theta: \text{Alg}(\mathcal{C}) \rightarrow \mathcal{C}$

HA3.1.3.14: Assume that the underlying ∞ -cat \mathcal{C} admits countable colimits, and \otimes preserves countable colimits. Then θ admits a left adjoint, which is given informally by the formula

$$\mathcal{C} \longmapsto \text{Sym}^* \mathcal{C} = \coprod_n \mathcal{C}^{\otimes n} / \Sigma_n$$

↑ " ∞ -categorical quotient by the symmetric group"

HA3.2: 1) If \mathcal{C} is complete (i.e. has small limits), then so is $\text{Alg}(\mathcal{C})$.

The forgetful functor θ detects (i.e. preserves and reflects) limits.

2) If \mathcal{C} is cocomplete, and \otimes is compatible with colimits, then $\text{Alg}(\mathcal{C})$

is cocomplete. The forgetful functor θ detects sifted colimits, but not general colimits.

(A simplicial set K is sifted if it is nonempty, and the diagonal map $K \rightarrow K \times K$ is cofinal.
Idea: sifted \supset filtered + quotient by equiv relations.)

3) If \mathcal{C} is presentable, and \otimes is compatible with colimits, then $\text{Alg}(\mathcal{C})$ is presentable.

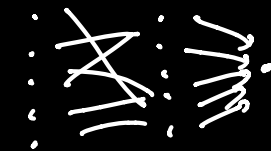
4) \exists a cocartesian symmetric monoidal structure on $\text{Alg}(\mathcal{C})$ st. the forgetful functor is monoidal. (Informally, the tensor product on \mathcal{C} induces a tensor product on comm alg. objects)

Def: **Associative operad** Assoc

It is a colored operad having a single object α

\forall finite set I , the set of operations $\text{Mul}_{\text{Assoc}}(\{\alpha\}_{i \in I}, \alpha)$ is the set of linear orderings on I .

Composition of linear orderings



\rightsquigarrow ∞ -operad $\text{Assoc}^{\otimes} \rightarrow \text{Fin}_*$ called **associative ∞ -operad**

\mathcal{C}^{\otimes} ∞ -operad equipped with a fibration $q: \mathcal{C}^{\otimes} \rightarrow \text{Assoc}^{\otimes}$

$\text{Alg}(\mathcal{C}) := \text{Alg}_{/\text{Assoc}}(\mathcal{C})$ ∞ -operad sections of q , called the **∞ -cat of associative algebra objects of \mathcal{C}** .

A **monoidal ∞ -cat** is a cocartesian fibration of ∞ -operads $\mathcal{C}^{\otimes} \rightarrow \text{Assoc}^{\otimes}$

HA4.1.1.14: Let $p: \mathcal{C}^{\otimes} \rightarrow \text{Fin}_*$ ∞ -operad. Then \mathcal{C}^{\otimes} is a symmetric

monoidal ∞ -cat if and only if the induced map $p': \text{Assoc}^{\otimes} \times_{\text{Fin}_*} \mathcal{C}^{\otimes} \rightarrow \text{Assoc}^{\otimes}$

is a monoidal ∞ -cat.

Def: **Left module operad** LM.

It is a colored operad having two elements n and m
 \forall objects $\{X_i\}_{i \in I}$ and Y of LM

$$\text{Mul}_{LM}(\{X_i\}_{i \in I}, Y) :=$$

{ If $Y = n$, $X_i = n \forall i$, then it is the collection of all linear orderings on I ;
If $Y = m$, and exactly one $X_i = m$, then it is the collection of linear orderings on I st. the last $X_i = m$.
In all other cases, it is \emptyset .

\rightsquigarrow ∞ -operad LM^{\otimes}

We have natural $\text{Assoc}^{\otimes} \longleftrightarrow LM^{\otimes} \xrightarrow{\text{fibration}} \text{Assoc}$

$\mathcal{C}^{\otimes} \rightarrow \text{Assoc}^{\otimes}$ fibration of ∞ -operads

∞ -category of left modules $L\text{Mod}(\mathcal{C}) := \text{Alg}_{LM/\text{Assoc}}(\mathcal{C})$

For $A \in \text{Alg}(\mathcal{C})$, we define ∞ -cat of left A -modules

$$L\text{Mod}_A(\mathcal{C}) := L\text{Mod}(\mathcal{C}) \times_{\text{Alg}(\mathcal{C})} \{A\}.$$

HA4.2: \mathcal{C} monoidal ∞ -cat, $A \in \text{Alg}(\mathcal{C})$, forgetful functor

$$\theta: L\text{Mod}_A(\mathcal{C}) \rightarrow \mathcal{C}.$$

(1) θ has a left adjoint, given by "free module".

(2) If \mathcal{C} is complete, then so is $L\text{Mod}_A(\mathcal{C})$, and θ detects limits.

(3) If \mathcal{C} is cocomplete, and \otimes preserves colimits, then $L\text{Mod}_A(\mathcal{C})$ is cocomplete, and θ detects colimits.

(4) If \mathcal{C} is presentable, and \otimes preserves colimits, then $L\text{Mod}_A(\mathcal{C})$ is presentable.

(5) Assume \mathcal{C} is symmetric monoidal, admits simplicial colimits, and \otimes preserves simplicial colimits. $A \in \text{Alg}(\mathcal{C})$.

Then $\text{Mod}_A(\mathcal{C}) := L\text{Mod}_A(\mathcal{C})$ is a symmetric monoidal ∞ -cat, and the forgetful functor θ is lax-monoidal (i.e. only a map of ∞ -operads).